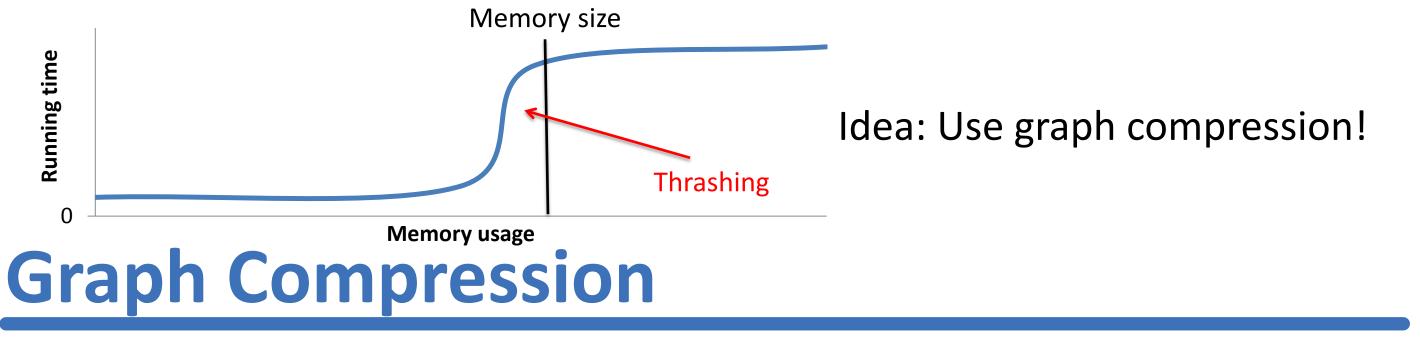
Ligra++: Processing Large Graphs Using Compression Julian Shun, Laxman Dhulipala, Guy Blelloch **Carnegie Mellon University**

Motivation

- Growth in graph data sizes (social networks, scientific computing, biology, etc.)
- Need to process graphs quickly
- What approach to use? Distributed memory, shared memory, disk-based
- Shared memory is the fastest, but limited by memory size

6

Cost of renting cloud machines increases with RAM size



Format: for each vertex, store differences between consecutive neighbors

Original

Ligra++

What about algorithm performance on compressed graphs? We implement graph compression and decoding techniques into the Ligra shared-memory graph processing framework

Ligra framework:

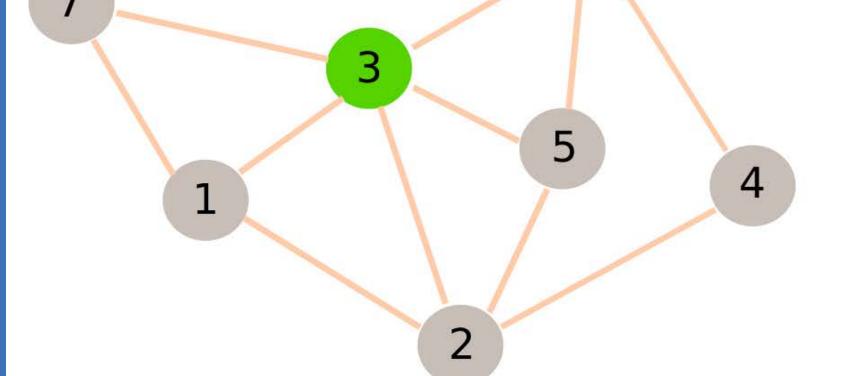
- represents a subset of vertices in a vertexSubset
- edgeMap: applies a function to the outgoing edges of a vertexSubset vertexMap: applies a function to the vertices in a vertexSubset

We modify the edgeMap function to decode each vertex's compressed edges on-the-fly

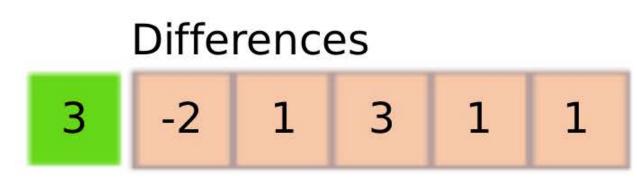
To allow for parallel decoding of high-degree vertices, we split the neighbors into chunks, compress each chunk separately, and decode each chunk in parallel

Encoding cost is amortized across all future computations on the graph

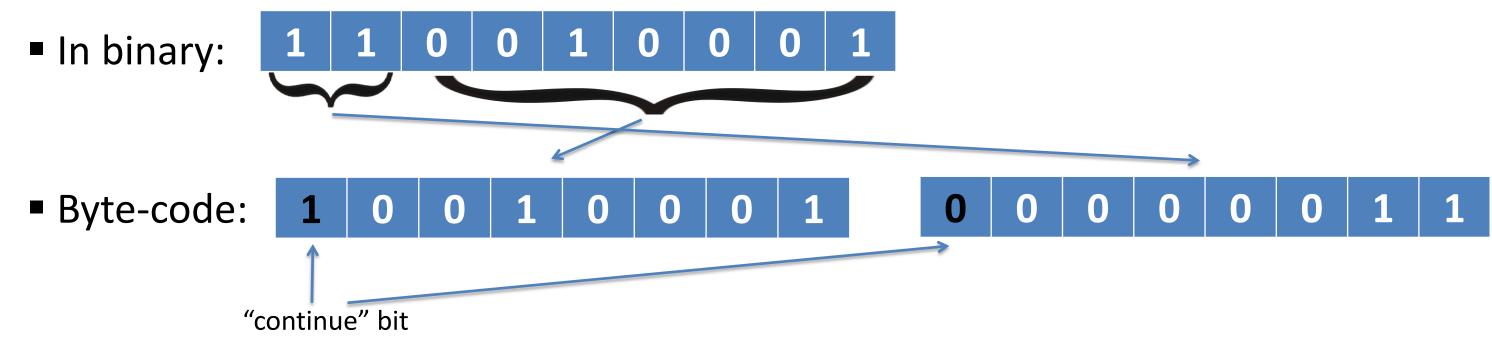
Performance







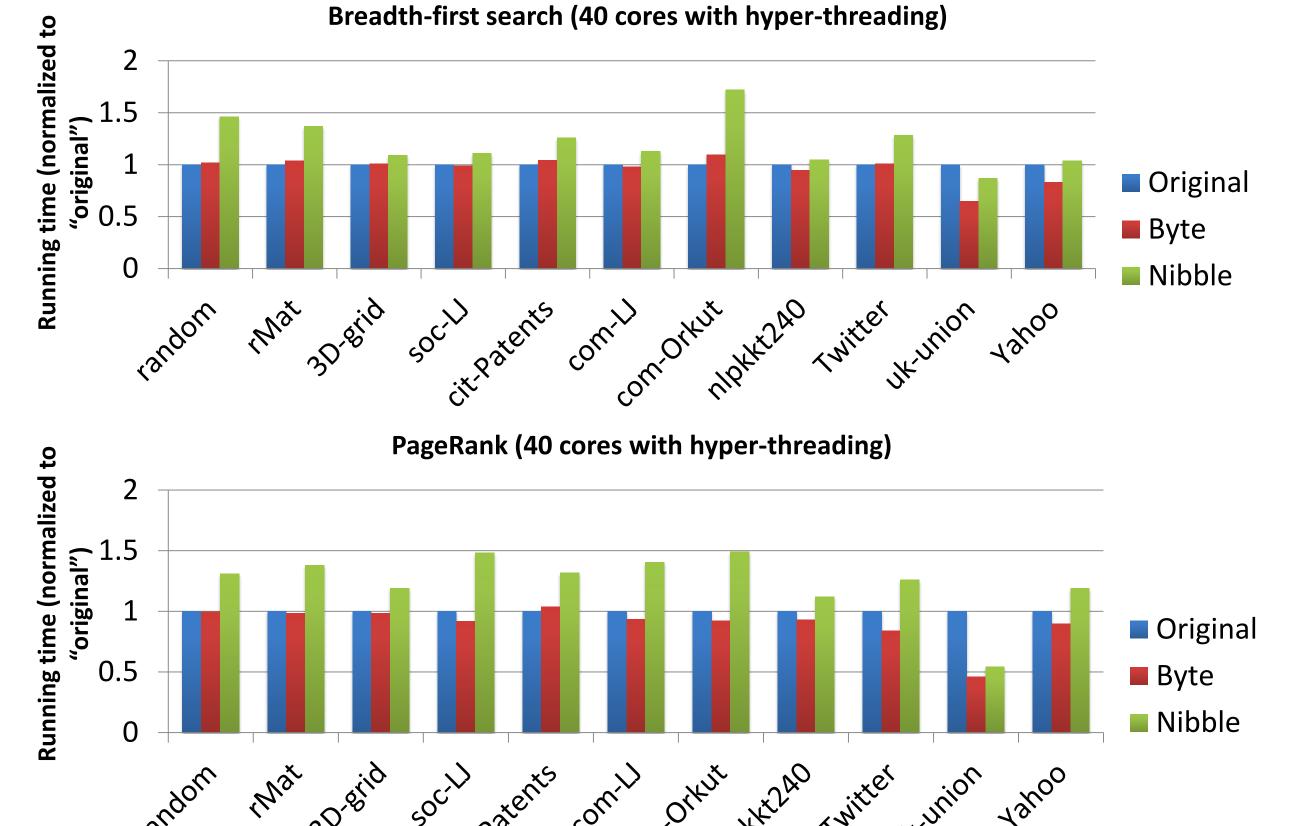
- Encode each difference using a k-bit code. Use k-1 bits for data, 1 bit as the "continue" bit
- We use 8-bit (byte) and 4-bit (nibble) codes
- Example: encode "401" using a byte-code



• Note: first difference can be negative, so the first code for it stores a "sign" bit. Decoding is just encoding "backwards" **Graph Reordering**

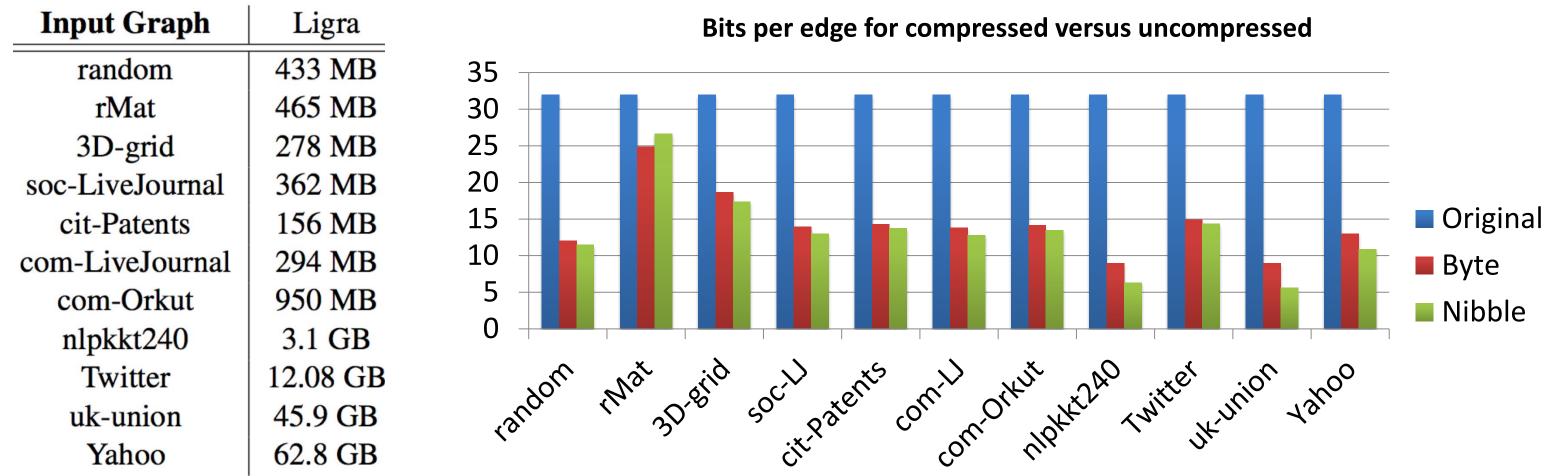
Can run graph reordering ("re-numbering") algorithms to improve locality

- Trade-offs: compressed versions have smaller memory footprint than uncompressed version, but requires time for decoding
- Performance of compressed versions much better in parallel than sequentially
- In parallel, memory bandwidth/contention is more of a bottleneck, and alleviates the cost of decoding!
- In parallel, byte code performance is competitive with uncompressed version



and compression (and also performance)

- Goal: have neighbors who have ID's close to own ID
- Various reordering algorithms: breadth-first search, depth-first search, hybrid BFS/DFS, METIS (based on finding graph separators), and our own separator-based algorithm
- Using best ordering, we get good compression for most graphs



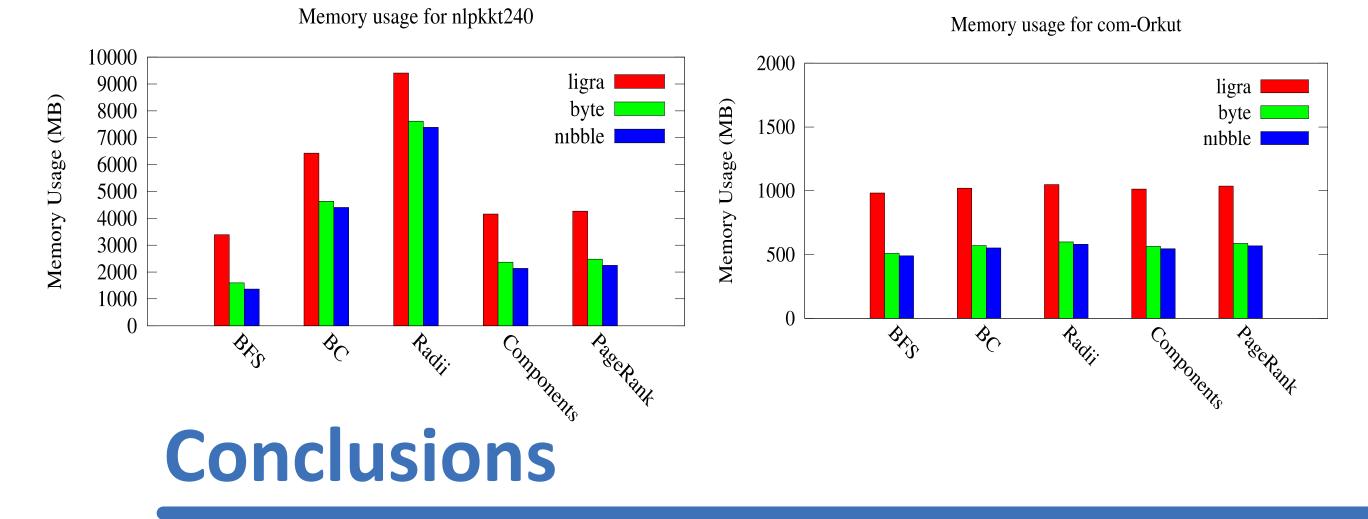
Running times on symmetrized Yahoo graph (1.4 billion vertices, 12.9 billion edges)

40-core Nehalem with hyper-threading	BFS	Betweenness Centrality	Radii	Connected Components		Bellman-Ford shortest paths
Original	4.66s	14s	24.5s	12s	8.27s	6.28s
Byte	3.87s	13.1s	23.5s	10.1s	7.47s	9.06s
Nibble	4.85s	18.6s	35.5s	15.7s	9.86s	13.7s

Mat 30-Brid social Patents conclusion or hit RADAO Twitter UN UKUNION random

Similar trends for other applications: betweenness centrality, radii estimation, connected components, and Bellman-Ford shortest paths

On 40 cores with hyper-threading, byte codes are between 1.5x slower and 2.7x faster



- With Ligra++, we can fit larger graphs than Ligra with the same amount of memory or the same graph with less memory while maintaining performance
- We are exploring techniques that reduce decoding cost to further improve the running time of Ligra++



