PACMAN: COORDINATED MEMORY CACHING FOR PARALLEL JOBS

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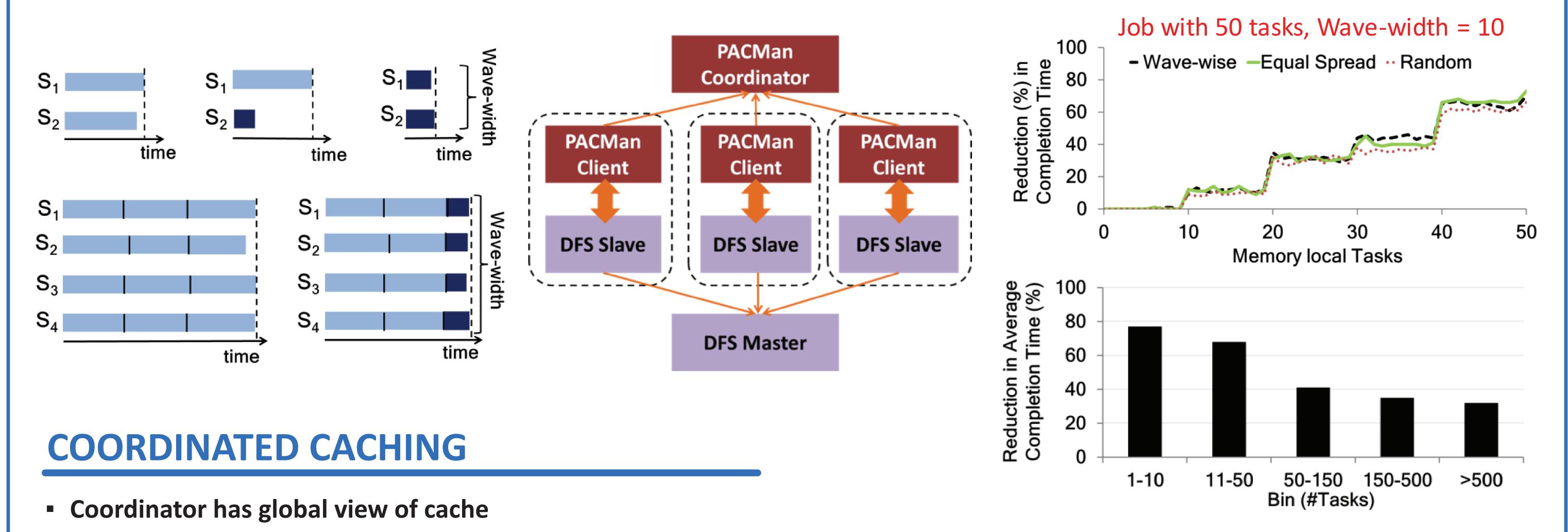
MOTIVATION

- MAnalytics jobs are parallel and process large amounts of data
- Machines have tens of gigabytes of memory
 - Falling memory prices
 - Median utilization of 19%
- **Heavy-tailed Input Sizes**

ALL-OR-NOTHING

- Jobs speed up when multiples of its wave-width are cached
 - Wave-width: #parallel executing tasks
 - Small single-waved jobs require 100% memory locality
- Cache hit-ratio insufficient; even MIN speeds up jobs by only 13%

- Elephant and mice jobs
- 92% of smallest job inputs can fit in memory
- Sticky policy: Focus replacements on incompletely cached waves



- Eviction and task placement
- Average Completion Time
 - LIFE: Evict from file with highest wave-width
 - Learn wave-width across multiple runs; file size correlates with wave-width
- Cluster Utilization
 - LFU-F: Evict from file with highest frequency
 - Overlap across map and reduce phases → sticky policy is important

EVALUATION

- Replayed Facebook and Bing workloads
- LIFE reduces average completion time by 53% and 51% in Facebook and Bing workloads
 - Small jobs see 77% improvement
- LFU-F improves cluster utilization by 47% and 53% in the Facebook and Bing workloads

PRE-FETCH AND PRE-REPLACE

- Oracle cache eviction and singly-accessed inputs
- Preparation Jobs: Large and multi-waved
- Pre-fetch for later waves of preparation jobs
- Evict inputs after multi-waved job ends
 - If singly-accessed, good!
 - If not, pre-fetch all but first wave
- Pre-replace with files of lowest wave-width

FUTURE WORK

- Proof of optimal cache eviction
- Hierarchical caching to include SSDs
- Details: HotOS 2011, NSDI 2012
- LIFE and LFU-F beat Belady's MIN despite lower cache hit-ratio
- Pre-fetch & Pre-replace → Ideal (87%) speedup
- Pre-replacement ~ Oracle cache eviction









